

Automatic Generation of Learning Assignments for Software Engineering Formalisms

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Learning Assignments

- ▶ Definition Seel [1981]

Learning assignments are “[...]selected and prepared learning objects with the aim to initiate, control, and organise learning processes.”

Software Engineering Formalism: Decision Tables

cf. Balzert [2009]

T	r_1	r_2	r_3
c_1	—	×	—
c_2	—	×	—
c_3	×	*	—
a_1	×	—	—
a_2	×	×	×

Decision Tables, Syntax

- ▶ Let C be a set of conditions and A be a set of actions, $C \cap A = \emptyset$
- ▶ A decision table T over C and A is a labelled $(m + k) \times n$ matrix

T	r_1	\dots	r_n
c_1	$v_{1,1}$	\dots	$v_{1,n}$
\vdots	\vdots	\ddots	\vdots
c_m	$v_{m,1}$	\dots	$v_{m,n}$
a_1	$w_{1,1}$	\dots	$w_{1,n}$
\vdots	\vdots	\ddots	\vdots
a_k	$w_{k,1}$	\dots	$w_{k,n}$

- ▶ where
 - ▶ $v_{1,1}, \dots, v_{m,n} \in \{\times, -, *\}$
 - ▶ $w_{1,1}, \dots, w_{k,n} \in \{\times, -\}$

Decision Tables, Semantics

Each rule $r \in \{r_1, \dots, r_n\}$ of table T

T	r_1	\dots	r_n
c_1	$v_{1,1}$	\dots	$v_{1,n}$
\vdots	\vdots	\ddots	\vdots
c_m	$v_{m,1}$	\dots	$v_{m,n}$
a_1	$w_{1,1}$	\dots	$w_{1,n}$
\vdots	\vdots	\ddots	\vdots
a_k	$w_{m,1}$	\dots	$w_{m,n}$

is assigned a propositional logical formula $\mathcal{F}(r)$ as follows:

Decision Tables, Semantics

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\vdots	\vdots	\ddots	\vdots
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is assigned a propositional logical formula $\mathcal{F}(r)$ as follows:

- ▶ Let (v_1, \dots, v_m) and (w_1, \dots, w_k) be premise and effect of r .
- ▶ Then

$$\mathcal{F}(r) := \underbrace{\bigwedge_{1 \leq i \leq m} F(v_i, c_i)}_{=: \mathcal{F}_{pre}(r)} \wedge \underbrace{\bigwedge_{1 \leq j \leq k} F(w_j, a_j)}_{=: \mathcal{F}_{eff}(r)}$$

where

$$F(v, x) = \begin{cases} x & \text{if } v = \times \\ \neg x & \text{if } v = - \\ true & \text{if } v = * \end{cases}$$

Decision Tables, Example

T	r_1	r_2	r_3
c_1	—	×	—
c_2	—	×	—
c_3	×	*	—
a_1	×	—	—
a_2	×	×	×

► $\mathcal{F}(r_1) = (\neg c_1 \wedge \neg c_2 \wedge c_3) \wedge (a_1 \wedge a_2)$

Decision Tables, Determinism

A decision table T is called deterministic if and only if the premises of all rules are pairwise disjoint, i.e. if

$$\forall r_1 \neq r_2 \in T \bullet \models \neg(\mathcal{F}_{pre}(r_1) \wedge \mathcal{F}_{pre}(r_2)).$$

Otherwise, T is called non-deterministic.

SE-Formalism: Decision Tables, Example

T	r_1	r_2	r_3
c_1	—	×	—
c_2	—	×	—
c_3	×	*	—
a_1	×	—	—
a_2	×	×	×

SE-Formalism: Decision Tables, Example

T	r_1	r_2	r_3
c_1	—	×	—
c_2	—	×	—
c_3	×	*	—
a_1	×	—	—
a_2	×	×	×

- ▶ T is deterministic, because no two rules can be enabled at the same time.

Decision Table Learning Assignments

T	r_1	r_2	r_3
c_1	×	×	—
c_2	×	—	*
c_3	—	×	*
a_1	×	—	—
a_2	—	×	—

1. Give the rule formulae for r_1, r_2, r_3 .
2. It is claimed that T is deterministic. Prove this claim.

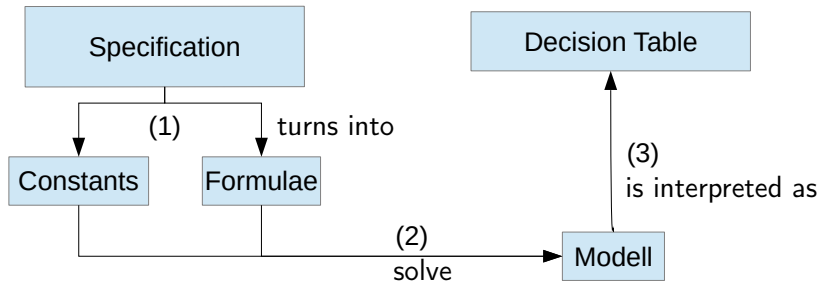
Decision Table Learning Assignments

T	r_1	r_2	r_3
c_1	×	×	—
c_2	×	—	*
c_3	—	×	*
a_1	×	—	—
a_2	—	×	—

1. Give the rule formulae for r_1, r_2, r_3 .
2. It is claimed that T is deterministic. Prove this claim.

- ▶ Creating learning assignments is hard.
 - ▶ Consistency constraints
 - ▶ Quality can vary
- ▶ Consistency constraints are formally and precisely defined
- ▶ Decision tables are mathematical objects

Generating Decision Tables, Overview



Generating Decision Tables, Consistency

T	r_1	r_2
c_1	$v_{1,1}$	$v_{1,2}$
c_2	$v_{2,1}$	$v_{2,2}$

Generating Decision Tables, Consistency

T	r_1	r_2
c_1	$v_{1,1}$	$v_{1,2}$
c_2	$v_{2,1}$	$v_{2,2}$

$$\forall r_1 \neq r_2 \in T \bullet \models \neg(\mathcal{F}_{pre}(r_1) \wedge \mathcal{F}_{pre}(r_2))$$

Generating Decision Tables, Consistency

T	r_1	r_2
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$$\forall r_1 \neq r_2 \in T \bullet \forall c_1, c_2 \bullet \neg(\mathcal{F}_{pre}(r_1) \wedge \mathcal{F}_{pre}(r_2))$$

Generating Decision Tables, Consistency

T	r_1	r_2
c_1	$v_{1,1}$	$v_{1,2}$
c_2	$v_{2,1}$	$v_{2,2}$

$$\forall r_1 \neq r_2 \in T \bullet \forall c_1, c_2 \bullet \neg((F(v_{1,1}, c_1) \wedge F(v_{2,1}, c_2)) \wedge (F(v_{1,2}, c_1) \wedge F(v_{2,2}, c_2)))$$

Generating Decision Tables, Consistency

T	r_1	r_2
c_1	$v_{1,1}$	$v_{1,2}$
c_2	$v_{2,1}$	$v_{2,2}$

$$\forall r_1 \neq r_2 \in T \bullet \forall c_1, c_2 \bullet \neg((G_{1,1} \wedge G_{2,1}) \wedge (G_{1,2} \wedge G_{2,2}))$$

with

$$G_{i,j} := (v_{i,j} = \times \wedge c_i) \vee (v_{i,j} = - \wedge \neg c_i) \vee (v_{i,j} = *)$$

Generating Decision Tables, Consistency

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with

$$G_{i,j} := (v_{i,j} = \times \wedge c_i) \vee (v_{i,j} = - \wedge \neg c_i) \vee (v_{i,j} = *)$$

satisfiable with $v_{1,1}, \dots, v_{2,2}$ free
results in some deterministic decision table

Generating Decision Tables, Quality

- ▶ each symbol should appear

$$\forall s \in \{\times, -, *\} \bullet (v_{1,1} = s) \vee (v_{1,2} = s) \vee (v_{2,1} = s) \vee (v_{2,2} = s)$$

- ▶ rules not the same

$$(v_{1,1} \neq v_{1,2}) \vee (v_{2,1} \neq v_{2,2})$$

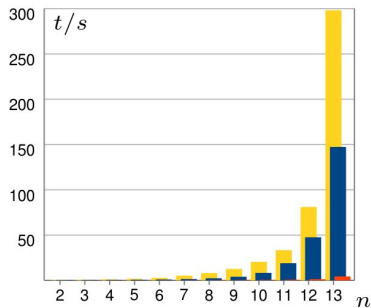
- ▶ counting *

$$n = (v_{1,1} = *) + (v_{1,2} = *) + (v_{2,1} = *) + (v_{2,2} = *)$$

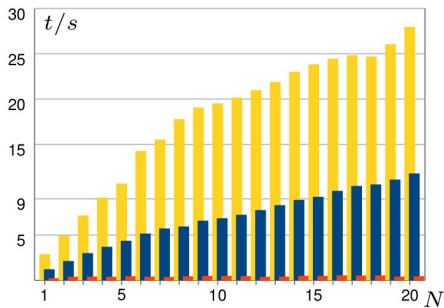
- ▶ ...

Results in some good decision table

Evaluation



(a) Generation of one decision table of size $(n + n) \times n$.



(b) Generation of a new bunch of 100 decision tables of size $(3 + 3) \times 3$.

Applications and Future Work

Possible Applications:

- ▶ check own learning assignment
- ▶ generate a learning assignment
- ▶ generate a different learning assignment
- ▶ for exams
- ▶ for exercise sheets
- ▶ for self-learning
- ▶ ...

Applications and Future Work

Possible Applications:

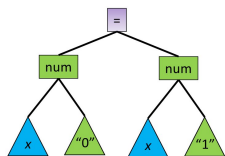
- ▶ check own learning assignment
- ▶ generate a learning assignment
- ▶ generate a different learning assignment
- ▶ for exams
- ▶ for exercise sheets
- ▶ for self-learning
- ▶ ...

Future Work:

- ▶ more decision table properties
- ▶ more quality formalizations
- ▶ other formalisms
- ▶ ...

Related Works

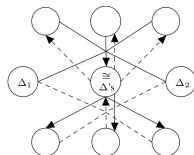
Shenoy, Aparanji, Sripradha, and Kumar [2016]



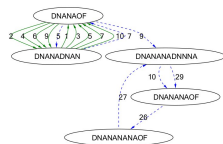
Singh, Gulwani, and Rajamani [2012]

$$\lim_{n \rightarrow \infty} \sum_{i=0}^n \frac{C_0 i^2 + C_1 i + C_2}{(C_3)^i} = \frac{C_4}{C_5}$$

Alvin, Gulwani, Majumdar, and Mukhopadhyay [2015]



Andersen, Gulwani, and Popovic [2013]



Fin

Thank you - Questions?

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